

Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application:

Listing of Claims:

1 Claim 1 (currently amended): A method of determining data flow for a channel having a  
2 plurality of subchannels in a multi-carrier system, comprising:

3 determining data flow for the channel ~~in terms of~~based on an input intensity  $\lambda_{in}$ ,  
4 ~~and a probability of having a frame having no or a correctable number of errors  $p$ , and a~~  
5 maximum number of transmissions  $k$  of the frame; and

6 adjusting channel performance in accordance with the data flow.

1 Claim 2 (currently amended): The method of claim 1 wherein said data flow is  
2 determined in accordance with the following relationships:  
3

4 
$$\lambda_{nac} = \lambda_{in} \frac{[1 - (1 - p)^k] (1 - p)}{p}$$

5 
$$\lambda_{ii} = \lambda_{in} \frac{1 - (1 - p)^k}{p},$$

6 
$$\lambda_{pout} = \lambda_{in} [1 - (1 - p)^k],$$

7 
$$\lambda_{rt} = \lambda_{in} \frac{(1 - p) [1 - (1 - p)^{k-1}]}{p}, \text{ and}$$

8 
$$\lambda_{nout} = \lambda_{in} (1 - p)^k,$$
  
9

10  $\lambda_{nac}$  represents a negative acknowledgement intensity,  ~~$k$  represents a maximum number of~~  
11 ~~transmissions~~,  $\lambda_{ti}$  represents a transmitter intensity,  $\lambda_{pout}$  represents an intensity of good  
12 and correctable frames;  $\lambda_{ri}$  represents a retransmission intensity;  $\lambda_{nout}$  represents an  
13 intensity of erroneous frames that are non-correctable after a maximum number of  
14 transmissions; and  
15 adjusting channel performance in accordance with the data flow.

1 Claim 3 (original): The method of claim 2 wherein the data flow is determined by  
2 applying said relationships to data flow in a downstream direction, and applying said  
3 relationships to data flow in an upstream direction.

1 Claim 4 (currently amended): A method of determining data flow for a channel having a  
2 plurality of subchannels in a multi-carrier system, comprising:  
3 determining an upstream data flow based on a maximum number transmissions of  
4 each frame;  
5 determining a downstream data flow based on the maximum number of  
6 transmissions of each frame; and  
7 superimposing the upstream data flow and the downstream data flow to determine  
8 a channel data flow.

1 Claim 5 (original): The method of claim 4 wherein the channel uses forward error  
2 correction.

Claim 6 (canceled)

Claim 7 (canceled)

Claim 8 (canceled)

Claim 9 (currently amended): A method of determining throughput in a multicarrier transmission system having a channel, comprising:  
generating a representation of the throughput of the channel in a first direction with respect to the throughput of the channel in a second direction, ~~The method of claim 8~~  
wherein the representation is generated in accordance with the following relationships:

$$\frac{M_d}{K_d} \left[ \frac{1}{m_d} + \frac{1-p_d}{p_d} \right] [1 - (1-p_d)^{k_d}] \Lambda_d + \frac{N_u}{K_u} \frac{1 - (1-p_u)^{k_u}}{p_u} \Lambda_u \leq V_u, \text{ and}$$

$$\frac{N_d}{K_d} \frac{1 - (1-p_d)^{k_d}}{p_d} \Lambda_d + \frac{M_u}{K_u} \left[ \frac{1}{m_u} + \frac{1-p_u}{p_u} \right] [1 - (1-p_u)^{k_u}] \Lambda_u \leq V_d,$$

wherein  $M_d$  represents a length of an acknowledgment frame in a downstream direction,  $K_d$  represents the length of an information field in the downstream direction,  $m_d$  represents a number of information frames between positive acknowledgment frames in the downstream direction,  $p_d$  represents a probability of an information frame being accepted in the downstream direction,  $k_d$  represents a maximum number of transmissions in the downstream direction,  $\Lambda_d$  represents a number of information bits per unit time in the downstream direction,  $N_d$  represents a total frame length in the downstream direction,  $M_u$  represents a length of an acknowledgment frame in an upstream direction,  $N_u$  represents a total frame length in the upstream direction,  $K_u$  represents the length of an information field in the upstream direction,  $m_u$  represents a number of information frames between positive acknowledgment frames in the upstream direction,  $p_u$  represents a probability of an information frame being accepted in the upstream direction,  $k_u$  represents a maximum number of transmissions in the upstream direction,  $\Lambda_u$  represents a number of information bits per unit time in the upstream direction,  $V_u$  represents a data rate in the upstream direction, and  $V_d$  represents a data rate in the downstream direction; and

determining the throughput of the channel in a first direction with respect to the  
throughput of the channel in a second direction using the representation.

Claim 10 (original): A method of determining throughput in a multicarrier transmission system, comprising:

determining the throughput of a channel in an upstream and downstream direction in accordance with the following relationships:

$$H_u = \max \Lambda_u = \min \left\{ \begin{aligned} &V_u / \left[ \frac{N_d}{K_d} \frac{1 - (1 - p_d)^{k_d}}{p_d} + \frac{M_u}{K_u} \frac{V_u}{V_d} \left( \frac{1}{m_u} + \frac{1 - p_u}{p_u} \right) (1 - (1 - p_u)^{k_u}) \right], \\ &V_u / \left[ \frac{M_d}{K_d} \frac{V_d}{V_u} \left( \frac{1}{m_d} + \frac{1 - p_d}{p_d} \right) (1 - (1 - p_d)^{k_d}) + \frac{N_u}{K_u} \frac{1 - (1 - p_u)^{k_u}}{p_u} \right] \end{aligned} \right\}, \text{ and}$$

$$H_d = \max \Lambda_d = \min \left\{ \begin{aligned} &V_d / \left[ \frac{N_d}{K_d} \frac{1 - (1 - p_d)^{k_d}}{p_d} + \frac{M_u}{K_u} \frac{V_u}{V_d} \left( \frac{1}{m_u} + \frac{1 - p_u}{p_u} \right) (1 - (1 - p_u)^{k_u}) \right], \\ &V_d / \left[ \frac{M_d}{K_d} \frac{V_d}{V_u} \left( \frac{1}{m_d} + \frac{1 - p_d}{p_d} \right) (1 - (1 - p_d)^{k_d}) + \frac{N_u}{K_u} \frac{1 - (1 - p_u)^{k_u}}{p_u} \right] \end{aligned} \right\},$$

wherein  $M_d$  represents a length of an acknowledgment frame in a downstream direction,  $K_d$  represents the length of an information field in the downstream direction,  $m_d$  represents a number of information frames between positive acknowledgment frames in the downstream direction,  $p_d$  represents a probability of an information frame being accepted in the downstream direction,  $k_d$  represents a maximum number of transmissions in the downstream direction,  $\Lambda_d$  represents a number of information bits per unit time in the downstream direction,  $N_d$  represents a total frame length in the downstream direction,  $M_u$  represents a length of an acknowledgment frame in an upstream direction,  $N_u$  represents a total frame length in the upstream direction,  $K_u$  represents the length of an information field in the upstream direction,  $m_u$  represents a number of information frames between positive acknowledgment frames in the upstream direction,  $p_u$  represents a probability of an information frame being accepted in the upstream direction,  $k_u$

represents a maximum number of transmissions in the upstream direction,  $\Lambda_u$  represents a number of information bits per unit time in the upstream direction,  $V_u$  represents a data rate in the upstream direction, and  $V_d$  represents a data rate in the downstream direction.

Claim 11 (currently amended): A method of increasing a bit load of a multicarrier system comprising a channel having a plurality of subchannels, comprising:

determining a bit load for at least one subchannel based on a target symbol error rate  $\varepsilon_s$ , a maximum number of symbol errors that can be corrected  $t$ , a number of symbols in an information field  $K$  of a frame, and a maximum number of transmissions  $k$  of the frame, and a number of bits per subchannel; and

selecting the maximum number of symbol errors  $t$ , the number of symbols in the information field  $K$  of the frame and the maximum number of transmissions  $k$  of the frame, such that a coding gain is increased.

Claim 12 (currently amended): The method of claim 11 wherein the coding gain is a function of an average number of transmissions for a frame.

Claim 13 (original): The method of claim 11 wherein the bit load is determined in accordance with the following relationships:

$$1 - \left( 1 - W(t, K, k) \varepsilon_s^{\frac{1}{(t+1)k}} \right)^{1/\alpha} \\ = \omega(b_i) (1 - 2^{-b_i/2}) \operatorname{erfc} \left( \sqrt{3 \cdot 10^{\gamma_i/10} / (2^{b_i+1} - 2)} \right) \left[ 2 - (1 - 2^{-b_i/2}) \operatorname{erfc} \left( \sqrt{3 \cdot 10^{\gamma_i/10} / (2^{b_i+1} - 2)} \right) \right]$$

wherein  $b_i$  represents a number of bits per subchannel,  $\gamma_i$  represents a signal-to-noise ratio at the  $i$ -th subchannel,  $\omega(b_i)$  represents an average fraction of erroneous bits in a  $b_i$ -sized erroneous quadrature-amplitude-modulation symbol,  $\varepsilon_s$  represents a target symbol error rate,  $\beta$  represents an effect of a descrambler, and  $\alpha$  represent a number of bits per code symbol; and

$$W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{\frac{k-1}{(t+1)k}}$$

C+R represents a number of redundant symbols in an error correction field, and the net coding gain  $G_n(t, K, k)$  is determined in accordance with the following relationship:

$$G_n(t, K, k) = \frac{K}{K+C+R} \frac{B_{DMT}(t, K, k)}{\nu} - \frac{K}{K+C} B_{DMT}(0, K, 1)$$

$\nu$  represents an average number of transmissions, and  $B_{DMT}(t, K, k)$  represents a number of bits in a discrete multitone symbol based on the values of t, K and k.

Claim 14 (original): The method of claim 13 wherein  $\omega(b_i)$  is approximated in accordance with the following relationship:

$$\omega(b_i) = \frac{4}{3 + 2b_i}$$

Claim 15 (original): The method of claim 13 wherein  $\varepsilon_s$  is determined in accordance with the following relationship:

$$\varepsilon_s = 1 - \left( 1 - \frac{\varepsilon}{\beta} \right)^\alpha,$$

and  $\varepsilon$  represents a target bit error rate,  $\alpha$  represents a length of a code symbol, and  $\beta$  represents the effect of a descrambler.

Claim 16 (original): The method of claim 13 wherein  $\omega(b_i)$  is determined in accordance with the following relationship:

$$\omega(b) = \frac{1}{b \cdot 2^b} \sum_{i=1}^{2^b} \sum_{j \neq i} \frac{d_H(a_i, a_j)}{\chi_i},$$

$b$  represents a number of bit positions of a quadrature-amplitude-modulation symbol,  $a_i$  represents a label for the  $i^{\text{th}}$  point of a constellation associated with a subchannel,  $a_j$  represents a label for the  $j^{\text{th}}$  point of a constellation associated with a subchannel, and  $\chi_i$  represents a coordination number of the  $a_i^{\text{th}}$  point,  $d_H(a_i, a_j)$  represents a Hamming distance between respective binary vectors associated with points  $a_i$  and  $a_j$ .

Claim 17 (original): The method of claim 11 further comprising:  
determining a total increase in the number of bits to be sent in a DMT symbol  
( $G_i(t, K, k)$ ) in accordance with the following relationship:

$$G_d(t, K, k) \equiv B_{DMT}(t, K, k) - B_{DMT}(0, K, 1).$$

Claim 18 (currently amended): A method of determining an uncoded bit error rate  $p_b$  based on a target symbol error rate  $\varepsilon_s$  and a maximum number of transmissions k of a frame, comprising:  
determining the uncoded bit error rate  $p_b$  based on a weighted series expansion of the target bit error rate  $\varepsilon_s$ , comprising weights  $W$  that are a function of a maximum number of symbol errors that can be corrected  $t$  and a number of symbols in an information field K of the frame; and  
selecting the maximum number of symbol errors  $t$ , the number of symbols in the information field K of the frame and the maximum number of transmissions k of the frame, such that the uncoded bit error rate  $p_b$  that produces a symbol error rate that is less than or equal to the target symbol error rate  $\varepsilon_s$  is largest.

1 Claim 19 (original): The method of claim 18 wherein said weighted series expansion to  
2 determine said uncoded bit error rate  $p_b$  comprises the following relationship:

3

$$4 \quad p_b = 1 - \left( 1 - W(t, K, k) \varepsilon_S^{\frac{1}{(t+1)k}} \right)^{1/\alpha}$$

5  
6 wherein

7

$$8 \quad W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{-\frac{k-1}{(t+1)k}},$$

9  
10 C+R represents a number of redundant symbols in an error correction field.

1 Claim 20 (currently amended): A method of selecting transmission parameters a  
2 multicarrier system having a channel comprising a plurality of subchannels, comprising:  
3 selecting a number (s) of discrete multi-tone symbols in a  
4 forward-error-correction frame, a number (z) of forward-error-correction control symbols  
5 in a discrete multitone symbol, and a maximum number of transmissions (k) of a frame,  
6 based on a signal-to-noise ratio and a number of subchannels associated with the  
7 signal-to-noise ratio; and

8 transmitting information in accordance with the selected number (s) of discrete  
9 multi-tone symbols, the number (z) of forward-error-correction control symbols in the  
10 discrete multitone symbol and the maximum number of transmissions (k) of the frame.

1 Claim 21 (original): The method of claim 20 wherein said selecting comprises selecting  
2 an adjustment value per subchannel based on the signal-to-noise ratio and the number of  
3 subchannels associated with the signal-to-noise ratio; and



4           adjusting a number of bits per subchannel for at least one subchannel in  
5           accordance with the adjustment value.

1       Claim 22 (original): The method of claim 20 wherein the signal-to-noise ratio is an  
2       average signal-to-noise ratio of the associated number of subchannels.

1       Claim 23 (currently amended): The method of claim 20 further comprising:  
2           storing, in a table, the number (s) of discrete multi-tone symbols in the  
3       forward-error-correction frame, the number (z) of forward-error-correction control  
4       symbols in the discrete multitone symbol associated with the signal-to-noise ratio, the  
5       maximum number of transmissions (k) of a frame, and the number of subchannels  
6       associated with the signal-to-noise ratio, for different values of s, z, signal-to-noise ratios  
7       and numbers of subchannels.

1       Claim 24 (original): The method of claim 23 wherein for each value of signal-to-noise  
2       ratio and number of bits per subchannel of the table, the associated values of s, z and k  
3       are also associated with an adjustment value that provides a maximal net coding gain  $g_n$ ,  
4       such that the associated values of s, z and k is selected from a subset of associated s, z  
5       and k values.

1       Claim 25 (original): A method of determining an optimum bit load b per subchannel in a  
2       multicarrier system with forward error correction, comprising:  
3           computing one or more values of a maximum number of symbol errors that can be  
4       corrected t, a number of symbols in the information field K and a maximum number of  
5       transmissions k to determine the optimum bit load per subchannel in accordance with the  
6       following relationship:

7  
8       
$$b = \lceil \gamma + \Phi(\gamma, t, K, k, \epsilon) \rceil / 10 \log 2$$
  
9

wherein

$$\Phi(\gamma, t, K, k, \varepsilon) = 10 \log \left\{ 10^{-\gamma/10} + \frac{3 \log e}{2 \log \left[ \frac{\alpha \langle \omega(b) \rangle \sqrt{8/\pi}}{W(t, K, k) (\alpha \varepsilon / \beta)^{\frac{1}{(t+1)k}}} \right] - \log \log \left[ \frac{\alpha \langle \omega(b) \rangle \sqrt{8/\pi}}{W(t, K, k) (\alpha \varepsilon / \beta)^{\frac{1}{(t+1)k}}} \right] + \log \left( \frac{\log e}{2} \right)} \right\}$$

$$W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{-\frac{k-1}{(t+1)k}},$$

$$\langle \omega(b) \rangle = \frac{1}{b_{\max}} \int_1^{b_{\max}} \omega(b) (1 - 2^{-b/2}) db$$

$\alpha$  represents a number of bits per symbol,  $\gamma$  represents a signal-to-noise ratio,  $\varepsilon$  represents a target symbol error rate,  $k$  represents a maximum number of transmissions,  $C+R$  represents a number of redundant symbols in an error correction field,  $b$  represents a number of bit positions of a quadrature-amplitude-modulation symbol,  $\omega(b)$  represents an average fraction of erroneous bits in an erroneous  $b$ -sized quadrature-amplitude-modulation symbol, and  $b_{\max}$  is a maximum number of bit positions of the quadrature-amplitude-modulation symbol per subchannel; and

selecting a bit load per subchannel in accordance with the maximum number of symbol errors that can be corrected  $t$ , a number of symbols in the information field  $K$  and the maximum number of transmissions  $k$ .

Claim 26 (original): A method for transmitting data in a multi-carrier system between a downstream station and an upstream station, coupled by a channel having a plurality of subchannels, comprising:

- transmitting an information frame from the upstream station;
- receiving the information frame at the downstream station;
- determining whether the information frame is non-correctable;

7           transmitting a negative acknowledgement when the information frame is  
8 non-correctable; and  
9           transmitting the information frame if the information frame has not be transmitted  
10 a predetermined number of times from the upstream station.

1       Claim 27 (original): The method of claim 26 wherein the predetermined number of times  
2 is determined in accordance with a measured signal-to-noise ratio value representing at  
3 least a subset of the subchannels of the channel, and forward error correction parameters.

1       Claim 28 (original): The method of claim 26 wherein the multi-carrier system is a  
2 discrete multi-tone system.

1       Claim 29 (original): The method of claim 26 wherein the discrete multi-tone system  
2 comprises the G-lite standard.

1       Claim 30 (original): The method of claim 26 wherein the discrete multi-tone system  
2 comprises the G.dmt standard.

1       Claim 31 (original): The method of claim 26 wherein the forward error correction  
2 parameters are Reed-Solomon forward error correction parameters.

Claim 32 (canceled)

1       Claim 33 (currently amended): ~~The apparatus of claim 32~~ An apparatus for determining  
2 throughput in a multicarrier transmission system having a channel, comprising:  
3       means for generating a representation of the throughput of the channel in a first  
4 direction with respect to the throughput of the channel in a second direction wherein the  
5 representation is generated in accordance with the following relationships:  
6

$$\frac{M_d}{K_d} \left[ \frac{1}{m_d} + \frac{1-p_d}{p_d} \right] \left[ 1 - (1-p_d)^{k_d} \right] \Lambda_d + \frac{N_u}{K_u} \frac{1 - (1-p_u)^{k_u}}{p_u} \Lambda_u \leq V_u, \text{ and}$$

$$\frac{N_d}{K_d} \frac{1 - (1-p_d)^{k_d}}{p_d} \Lambda_d + \frac{M_u}{K_u} \left[ \frac{1}{m_u} + \frac{1-p_u}{p_u} \right] \left[ 1 - (1-p_u)^{k_u} \right] \Lambda_u \leq V_d,$$

wherein  $M_d$  represents a length of an acknowledgment frame in a downstream direction,  $K_d$  represents the length of an information field in the downstream direction,  $m_d$  represents a number of information frames between positive acknowledgment frames in the downstream direction,  $p_d$  represents a probability of an information frame being accepted in the downstream direction,  $k_d$  represents a maximum number of transmissions in the downstream direction,  $\Lambda_d$  represents a number of information bits per unit time in the downstream direction,  $N_d$  represents a total frame length in the downstream direction,  $M_u$  represents a length of an acknowledgment frame in an upstream direction,  $N_u$  represents a total frame length in the upstream direction,  $K_u$  represents the length of an information field in the upstream direction,  $m_u$  represents a number of information frames between positive acknowledgment frames in the upstream direction,  $p_u$  represents a probability of an information frame being accepted in the upstream direction,  $k_u$  represents a maximum number of transmissions in the upstream direction,  $\Lambda_u$  represents a number of information bits per unit time in the upstream direction,  $V_u$  represents a data rate in the upstream direction, and  $V_d$  represents a data rate in the downstream direction; and

means for determining the throughput of the channel in a first direction with respect to the throughput of the channel in a second direction using the representation.

Claim 34 (original): An apparatus for determining throughput in a multicarrier transmission system, comprising:

means for determining the throughput of a channel in an upstream and downstream direction in accordance with the following relationships:

$$H_u = \max \Lambda_u = \min \left\{ \begin{array}{l} V_u / \left[ \frac{N_d}{K_d} \frac{1 - (1 - p_d)^{k_d}}{p_d} + \frac{M_u}{K_u} \frac{V_u}{V_d} \left( \frac{1}{m_u} + \frac{1 - p_u}{p_u} \right) (1 - (1 - p_u)^{k_u}) \right], \\ V_u / \left[ \frac{M_d}{K_d} \frac{V_d}{V_u} \left( \frac{1}{m_d} + \frac{1 - p_d}{p_d} \right) (1 - (1 - p_d)^{k_d}) + \frac{N_u}{K_u} \frac{1 - (1 - p_u)^{k_u}}{p_u} \right] \end{array} \right\}, \text{ and}$$

$$H_d = \max \Lambda_d = \min \left\{ \begin{array}{l} V_d / \left[ \frac{N_d}{K_d} \frac{1 - (1 - p_d)^{k_d}}{p_d} + \frac{M_u}{K_u} \frac{V_u}{V_d} \left( \frac{1}{m_u} + \frac{1 - p_u}{p_u} \right) (1 - (1 - p_u)^{k_u}) \right], \\ V_d / \left[ \frac{M_d}{K_d} \frac{V_d}{V_u} \left( \frac{1}{m_d} + \frac{1 - p_d}{p_d} \right) (1 - (1 - p_d)^{k_d}) + \frac{N_u}{K_u} \frac{1 - (1 - p_u)^{k_u}}{p_u} \right] \end{array} \right\},$$

wherein  $M_d$  represents a length of an acknowledgment frame in a downstream direction,  $K_d$  represents the length of an information field in the downstream direction,  $m_d$  represents a number of information frames between positive acknowledgment frames in the downstream direction,  $p_d$  represents a probability of an information frame being accepted in the downstream direction,  $k_d$  represents a maximum number of transmissions in the downstream direction,  $\Lambda_d$  represents a number of information bits per unit time in the downstream direction,  $N_d$  represents a total frame length in the downstream direction,  $M_u$  represents a length of an acknowledgment frame in an upstream direction,  $N_u$  represents a total frame length in the upstream direction,  $K_u$  represents the length of an information field in the upstream direction,  $m_u$  represents a number of information frames between positive acknowledgment frames in the upstream direction,  $p_u$  represents a probability of an information frame being accepted in the upstream direction,  $k_u$  represents a maximum number of transmissions in the upstream direction,  $\Lambda_u$  represents a number of information bits per unit time in the upstream direction,  $V_u$  represents a data rate in the upstream direction, and  $V_d$  represents a data rate in the downstream direction.

Claim 35 (currently amended): An apparatus for increasing a bit load of a multicarrier system comprising a channel having a plurality of subchannels, comprising:

means for determining a bit load for at least one subchannel based on a target symbol error rate  $\varepsilon_s$ , a maximum number of symbol errors that can be corrected  $t$ , a number of symbols in an information field  $K$  of a frame, and a maximum number of transmissions  $k$  of the frame, and a number of bits per subchannel; and

means for selecting the maximum number of symbol errors  $t$ , the number of symbols in the information field  $K$  of the frame and the maximum number of transmissions  $k$  of the frame, such that a coding gain is increased.

Claim 36 (original): The apparatus of claim 35 wherein the coding gain is a function of an average number of transmissions for a frame.

Claim 37 (original): The apparatus of claim 35 wherein the bit load is determined in accordance with the following relationships:

$$1 - \left( 1 - W(t, K, k) \varepsilon_s^{\frac{1}{(t+1)k}} \right)^{1/\alpha} \\ = \omega(b_i) (1 - 2^{-b_i/2}) \operatorname{erfc} \left( \sqrt{3 \cdot 10^{\gamma_i/10} / (2^{b_i+1} - 2)} \right) \left[ 2 - (1 - 2^{-b_i/2}) \operatorname{erfc} \left( \sqrt{3 \cdot 10^{\gamma_i/10} / (2^{b_i+1} - 2)} \right) \right]$$

wherein  $b_i$  represents a number of bits per subchannel,  $\gamma_i$  represents a signal-to-noise ratio at the  $i$ -th subchannel,  $\omega(b_i)$  represents an average fraction of erroneous bits in a  $b_i$ -sized erroneous quadrature-amplitude-modulation symbol,  $\varepsilon_s$  represents a target symbol error rate,  $\beta$  represents an effect of a descrambler, and  $\alpha$  represent a number of bits per code symbol; and

$$W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{-\frac{k-1}{(t+1)k}}$$

C+R represents a number of redundant symbols in an error correction field, and the net coding gain  $G_n(t, K, k)$  is determined in accordance with the following relationship:

$$G_n(t, K, k) = \frac{K}{K + C + R} \frac{B_{DMT}(t, K, k)}{\nu} - \frac{K}{K + C} B_{DMT}(0, K, 1)$$

$\nu$  represents an average number of transmissions, and  $B_{DMT}(t, K, k)$  represents a number of bits in a discrete multitone symbol based on the values of t, K and k.

Claim 38 (original): The apparatus of claim 37 wherein  $\omega(b_i)$  is approximated in accordance with the following relationship:

$$\omega(b_i) = \frac{4}{3 + 2b_i}$$

Claim 39 (original): The apparatus of claim 37 wherein  $\varepsilon_s$  is determined in accordance with the following relationship:

$$\varepsilon_s = 1 - \left(1 - \frac{\varepsilon}{\beta}\right)^\alpha,$$

and  $\varepsilon$  represents a target bit error rate,  $\alpha$  represents a length of a code symbol, and  $\beta$  represents the effect of a descrambler.

Claim 40 (original): The apparatus of claim 37 wherein  $\omega(b_i)$  is determined in accordance with the following relationship:

$$\omega(b) = \frac{1}{b \cdot 2^b} \sum_{i=1}^{2^b} \sum_{j \neq i}^{x_i} \frac{d_H(a_i, a_j)}{\chi_i},$$

5  
6  $b$  represents a number of bit positions of a quadrature-amplitude-modulation symbol,  $a_i$   
7 represents a label for the  $i^{\text{th}}$  point of a constellation associated with a subchannel,  $a_j$   
8 represents a label for the  $j^{\text{th}}$  point of a constellation associated with a subchannel, and  $\chi_i$   
9 represents a coordination number of the  $a_i^{\text{th}}$  point,  $d_H(a_i, a_j)$  represents a Hamming  
10 distance between respective binary vectors associated with points  $a_i$  and  $a_j$ .

1 Claim 41 (original): The apparatus of claim 35 further comprising:  
2 means for determining a total increase in the number of bits to be sent in a DMT  
3 symbol ( $G_i(t, K, k)$ ) in accordance with the following relationship:  
4

$$5 \quad G_i(t, K, k) \equiv B_{DMT}(t, K, k) - B_{DMT}(0, K, 1).$$

1 Claim 42 (currently amended): An apparatus for determining an uncoded bit error rate  $p_b$   
2 based on a target symbol error rate  $\varepsilon_s$  and a maximum number of transmissions  $k$  of a frame,  
3 comprising:

4 means for determining the uncoded bit error rate  $p_b$  based on a weighted series  
5 expansion of the target bit error rate  $\varepsilon_s$ , comprising weights  $W$  that are a function of a  
6 maximum number of symbol errors that can be corrected  $t$  and a number of symbols in an  
7 information field  $K$  of the frame; and

8 means for selecting the maximum number of symbol errors  $t$ , the number of symbols  
9 in the information field  $K$  of the frame and the maximum number of transmissions  $k$  of the  
10 frame, such that the uncoded bit error rate  $p_b$  that produces a symbol error rate that is less  
11 than or equal to the target symbol error rate  $\varepsilon_s$  is largest.

1 Claim 43 (original): The apparatus of claim 42 wherein said weighted series expansion to  
2 determine said uncoded bit error rate  $p_b$  comprises the following relationship:  
3



$$p_b = 1 - \left( 1 - W(t, K, k) \varepsilon_s^{\frac{1}{(t+1)k}} \right)^{1/\alpha}$$

wherein

$$W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{-\frac{k-1}{(t+1)k}}$$

C+R represents a number of redundant symbols in an error correction field.

Claim 44 (original): An apparatus for selecting transmission parameters a multicarrier system having a channel comprising a plurality of subchannels, comprising:

means for selecting a number (s) of discrete multi-tone symbols in a forward-error-correction frame, a number (z) of forward-error-correction control symbols in a discrete multitone symbol, and a maximum number of transmissions (k), based on a signal-to-noise ratio and a number of subchannels associated with the signal-to-noise ratio; and

means for transmitting information in accordance with the selected number (s) of discrete multi-tone symbols, the number (z) of forward-error-correction control symbols in the discrete multitone symbol and the maximum number of transmissions (k).

Claim 45 (original): The apparatus of claim 44 wherein said means for selecting comprises selecting an adjustment value per subchannel based on the signal-to-noise ratio and the number of subchannels associated with the signal-to-noise ratio; and

means for adjusting a number of bits per subchannel for at least one subchannel in accordance with the adjustment value.

1 Claim 46 (original): The apparatus of claim 44 wherein the signal-to-noise ratio is an  
2 average signal-to-noise ratio of the associated number of subchannels.

1 Claim 47 (original): The apparatus of claim 44 further comprising:  
2 means for storing, in a table, the number (s) of discrete multi-tone symbols in the  
3 forward-error-correction frame, the number (z) of forward-error-correction control  
4 symbols in the discrete multitone symbol associated with the signal-to-noise ratio, the  
5 maximum number of transmissions (k) and the number of subchannels associated with  
6 the signal-to-noise ratio, for different values of s, z, signal-to-noise ratios and numbers of  
7 subchannels.

1 Claim 48 (original): The apparatus of claim 47 wherein for each value of signal-to-noise  
2 ratio and number of bits per subchannel of the table, the associated values of s, z and k  
3 are also associated with an adjustment value that provides a maximal net coding gain  $g_n$ ,  
4 such that the associated values of s, z and k is selected from a subset of associated s, z  
5 and k values.

1 Claim 49 (original): An apparatus for determining an optimum bit load b per subchannel in a  
2 multicarrier system with forward error correction, comprising:  
3 means for computing one or more values of a maximum number of symbol errors that  
4 can be corrected  $t$ , a number of symbols in the information field  $K$  and a maximum  
5 number of transmissions  $k$  to determine the optimum bit load per subchannel in  
6 accordance with the following relationship:  
7

$$8 \quad b = \lceil \gamma + \Phi(\gamma, t, K, k, \varepsilon) \rceil / 10 \log 2$$

9  
10 wherein

$$\Phi(\gamma, t, K, k, \varepsilon) = 10 \log \left\{ 10^{-\gamma/10} + \frac{3 \log e}{2 \log \left[ \frac{\alpha \langle \omega(b) \rangle \sqrt{8/\pi}}{W(t, K, k) (\alpha \varepsilon / \beta)^{1/(t+1)k}} \right] - \log \log \left[ \frac{\alpha \langle \omega(b) \rangle \sqrt{8/\pi}}{W(t, K, k) (\alpha \varepsilon / \beta)^{1/(t+1)k}} \right] + \log \left( \frac{\log e}{2} \right)} \right\}$$

$$W(t, K, k) = \left[ \binom{K+C+R-1}{t} \right]^{-\frac{1}{(t+1)k}} \left[ \binom{K+C+R}{t+1} \right]^{-\frac{k-1}{(t+1)k}},$$

$$\langle \omega(b) \rangle = \frac{1}{b_{\max}} \int_1^{b_{\max}} \omega(b) (1 - 2^{-b/2}) db$$

$\alpha$  represents a number of bits per symbol,  $\gamma$  represents a signal-to-noise ratio,  $\varepsilon$  represents a target symbol error rate,  $k$  represents a maximum number of transmissions,  $C+R$  represents a number of redundant symbols in an error correction field,  $b$  represents a number of bit positions of a quadrature-amplitude-modulation symbol,  $\omega(b)$  represents an average fraction of erroneous bits in an erroneous  $b$ -sized quadrature-amplitude-modulation symbol, and  $b_{\max}$  is a maximum number of bit positions of the quadrature-amplitude-modulation symbol per subchannel; and  
means for selecting a bit load per subchannel in accordance with the maximum number of symbol errors that can be corrected  $t$ , a number of symbols in the information field  $K$  and the maximum number of transmissions  $k$ .

Claim 50 (original): A method for transmitting data in a multi-carrier system between a downstream station and an upstream station, coupled by a channel having a plurality of subchannels, comprising:

- a transmitter to transmit an information frame from the upstream station;
- a receiver to receive the information frame at the downstream station, the receiver to determine whether the information frame is non-correctable, and transmit a negative acknowledgement when the information frame is non-correctable;

8            wherein the transmitter, in response to the negative acknowledgment, transmits  
9            the information frame if the information frame has not be transmitted a predetermined  
10           number of times from the upstream station.

1           Claim 51 (original): The apparatus of claim 50 wherein the predetermined number of  
2           times is determined in accordance with a measured signal-to-noise ratio value  
3           representing at least a subset of the subchannels of the channel, and forward error  
4           correction parameters.

1           Claim 52 (original): The apparatus of claim 50 wherein the multi-carrier system is a  
2           discrete multi-tone system.

1           Claim 53 (original): The apparatus of claim 50 wherein the discrete multi-tone system  
2           comprises the G-lite standard.

1           Claim 54 (original): The apparatus of claim 50 wherein the discrete multi-tone system  
2           comprises the G.dmt standard.

1           Claim 55 (original): The apparatus of claim 50 wherein the forward error correction  
2           parameters are Reed-Solomon forward error correction parameters.